

**Problem 2.1 Min-Entropy in the i.i.d. limit**

Let us introduce the “smoothed” min-entropy of a random variable  $X$  over  $\mathcal{X}$  as

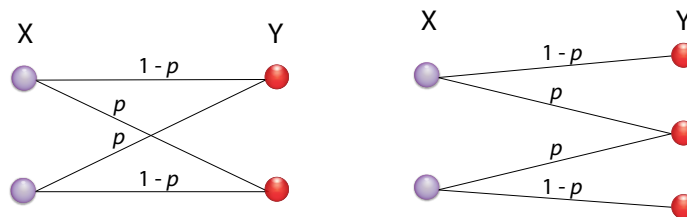
$$H_{\min}^\epsilon(X) = \max_{Q_X} \min_{x \in \mathcal{X}} h_Q(x), \tag{1}$$

$\delta(Q_X, P_X) < \epsilon$

where  $h_Q(x) = -\log Q_X(x)$  and the maximum is over all probability distributions  $Q_X$  that are  $\epsilon$ -close to  $P_X$ . Further, we define an i.i.d. random variable  $\vec{X} = \{X_1, X_2, \dots, X_n\}$  on  $\mathcal{X}^{\times n}$  with  $P_{\vec{X}}(\vec{x}) = \prod_{i=1}^n P_X(x_i)$ . Use the weak law of large numbers to show that the “smoothed” min-entropy converges to the Shannon entropy  $H(X)$  in the i.i.d. limit:

$$\lim_{\epsilon \rightarrow 0} \lim_{n \rightarrow \infty} \frac{1}{n} H_{\min}^\epsilon(\vec{X}) \geq H(X). \tag{2}$$

**Problem 2.2 Channel capacity**



(a) Binary Symmetric Channel      (b) Symmetric Erasure Channel

The asymptotic channel capacity is given by

$$C = \max_{P_X} I(X : Y).$$

- a) Calculate the asymptotic capacities of the two channels depicted above.
- b) Can we transmit a message error-free and with a finite amount of channel uses?
- \*c) Show that feedback does not improve the channel capacity.